

## §4 推理验证

### §4.1 一阶迁移系统的推理

例子: 互斥协议

$T$	$a = s_0$	$\longrightarrow$	$(y, t, a) := (1, 1, s_1)$
	$a = s_1 \wedge (x = 0 \vee t = 0)$	$\longrightarrow$	$(a) := (s_2)$
	$a = s_2$	$\longrightarrow$	$(y, a) := (0, s_3)$
	$a = s_3$	$\longrightarrow$	$(y, t, a) := (1, 1, s_1)$
	$b = t_0$	$\longrightarrow$	$(x, t, b) := (1, 0, t_1)$
	$b = t_1 \wedge (y = 0 \vee t = 1)$	$\longrightarrow$	$(b) := (t_2)$
	$b = t_2$	$\longrightarrow$	$(x, b) := (0, t_3)$
	$b = t_3$	$\longrightarrow$	$(x, t, b) := (1, 0, t_1)$
$\Theta$	$(a = s_0 \wedge b = t_0 \wedge x = 0 \wedge y = 0 \wedge t = 0)$		

安全性质: 互斥协议满足以下性质:

$$(T, \Theta) \models_I \Box(a = s_1 \wedge b \neq t_1 \wedge b \neq t_2 \rightarrow a = s_2 R b \neq t_2)$$

主要适用规则

主要适用规则如下

$$\frac{\begin{array}{l} \zeta \Rightarrow \varphi' \\ \{\varphi' \wedge \neg\psi\}T\{\varphi'\} \\ \varphi' \Rightarrow \varphi \end{array}}{\zeta \Rightarrow \psi R \varphi}$$

尝试 1

令

$$\begin{aligned} \zeta &\equiv (a = s_1 \wedge b \neq t_1 \wedge b \neq t_2) \\ \psi &\equiv (a = s_2) \\ \varphi &\equiv (b \neq t_2) \\ \varphi' &\equiv \varphi \end{aligned}$$

需要  $\zeta \Rightarrow \varphi$  和  $\forall t \in T, \{\varphi' \wedge \neg\psi\}t\{\varphi'\}$ 。

试证:

$$\{\varphi' \wedge \neg\psi\}t\{\varphi'\}.$$

需要

$$(b \neq t_2) \wedge (a \neq s_2) \wedge b = t_1 \wedge (y = 0 \vee t = 1) \rightarrow (t_2 \neq t_2)$$

证明失败。

## 尝试 2

令

$$\zeta \equiv (a = s_1 \wedge b \neq t_1 \wedge b \neq t_2)$$

$$\psi \equiv (a = s_2)$$

$$\varphi \equiv (b \neq t_2)$$

$$\varphi' \equiv (a = s_1 \wedge (b = t_0 \vee b = t_3 \vee (b = t_1 \wedge x = 1 \wedge t = 0))) \vee (a = s_2 \wedge b \neq t_2)$$

试证:

$$\{\varphi' \wedge \neg\psi\} t_6 \{\varphi'\}.$$

需要

$$((a = s_1 \wedge (b = t_0 \vee b = t_3 \vee (b = t_1 \wedge x = 1 \wedge t = 0))) \vee (a = s_2 \wedge b \neq t_2)) \wedge$$

$$(b \neq t_2) \wedge (a \neq s_2) \wedge b = t_1 \wedge (y = 0 \vee t = 1)$$

$$\rightarrow (a = s_1 \wedge (t_2 = t_0 \vee t_2 = t_3 \vee (t_2 = t_1 \wedge x = 1 \wedge t = 0))) \vee (a = s_2 \wedge t_2 \neq t_2)$$

$$((a = s_1 \wedge (b = t_0 \vee b = t_3 \vee (b = t_1 \wedge x = 1 \wedge t = 0))) \wedge$$

$$(b \neq t_2) \wedge (a \neq s_2) \wedge b = t_1 \wedge (y = 0 \vee t = 1)$$

$$\rightarrow (a = s_1 \wedge (t_2 = t_0 \vee t_2 = t_3 \vee (t_2 = t_1 \wedge x = 1 \wedge t = 0))) \vee (a = s_2 \wedge t_2 \neq t_2)$$

$$((a = s_1 \wedge ((b = t_1 \wedge x = 1 \wedge t = 0))) \wedge (y = 0))$$

$$\rightarrow (a = s_1 \wedge (t_2 = t_0 \vee t_2 = t_3 \vee (t_2 = t_1 \wedge x = 1 \wedge t = 0))) \vee (a = s_2 \wedge t_2 \neq t_2)$$

证明失败。

## 尝试 3

令

$$\zeta \equiv (a = s_1 \wedge b \neq t_1 \wedge b \neq t_2)$$

$$\psi \equiv (a = s_2)$$

$$\varphi \equiv (b \neq t_2)$$

$$\varphi' \equiv (a = s_1 \wedge (b = t_0 \vee b = t_3 \vee (b = t_1 \wedge x = 1 \wedge t = 0 \wedge y = 1))) \vee (a = s_2 \wedge b \neq t_2)$$

试证:

$$\{\varphi' \wedge \neg\psi\} t_6 \{\varphi'\}.$$

需要

$$((a = s_1 \wedge (b = t_0 \vee b = t_3 \vee (b = t_1 \wedge x = 1 \wedge t = 0 \wedge y = 1))) \vee (a = s_2 \wedge b \neq t_2)) \wedge$$

$$(b \neq t_2) \wedge (a \neq s_2) \wedge b = t_1 \wedge (y = 0 \vee t = 1)$$

$$\rightarrow (a = s_1 \wedge (t_2 = t_0 \vee t_2 = t_3 \vee (t_2 = t_1 \wedge x = 1 \wedge t = 0 \wedge y = 1))) \vee (a = s_2 \wedge t_2 \neq t_2)$$

$$((a = s_1 \wedge (b = t_0 \vee b = t_3 \vee (b = t_1 \wedge x = 1 \wedge t = 0 \wedge y = 1))) \wedge$$

$$(b \neq t_2) \wedge (a \neq s_2) \wedge b = t_1 \wedge (y = 0 \vee t = 1)$$

$$\rightarrow (a = s_1 \wedge (t_2 = t_0 \vee t_2 = t_3 \vee (t_2 = t_1 \wedge x = 1 \wedge t = 0 \wedge y = 1))) \vee (a = s_2 \wedge t_2 \neq t_2)$$

因此  $a = s_1 \wedge b \neq t_1 \wedge b \neq t_2 \Rightarrow a = s_2 \wedge b \neq t_2$ ,

即  $\square(a = s_1 \wedge b \neq t_1 \wedge b \neq t_2 \rightarrow a = s_2 \wedge b \neq t_2)$ 。

分两种情况来证明, 分析的时候可能方便些。

我们有  $\zeta \rightarrow \zeta_1 \vee \zeta_2$ ,

其中  $\zeta_1 = (a = s_1 \wedge b = t_0)$ ,  $\zeta_2 = (a = s_1 \wedge b = t_3)$ ,

令

$$\varphi_1 = (a = s_1 \wedge (b = t_0 \vee (b = t_1 \wedge x = 1 \wedge t = 0 \wedge y = 1))) \vee (a = s_2 \wedge b \neq t_2)$$

则  $\zeta_1 \Rightarrow \psi R \varphi_1$  .

因此  $\zeta_1 \Rightarrow a = s_2 R b \neq t_2$  .

同理  $\zeta_2 \Rightarrow a = s_2 R b \neq t_2$  .

因此  $\zeta \Rightarrow a = s_2 R b \neq t_2$  .

...

**活性性质:** 互斥协议满足以下性质:

$$(T, \Theta) \models_I \Box (a = s_1 \rightarrow \Diamond a = s_2)$$

**主要适用规则**

主要适用规则如下

$$\begin{array}{l} \varphi \Rightarrow (\psi \vee \zeta) \\ \zeta \Rightarrow (w_x^e \wedge (\psi \vee COND_T)) \\ \frac{\{\zeta \wedge e = v\} T \{\psi \vee (\zeta \wedge e \sqsubset v)\}}{\varphi \Rightarrow \Diamond \psi} \end{array}$$

**尝试 1**

选择  $f$  满足

$$\begin{array}{llll} I(f(t_0, 0)) = 1 & I(f(t_1, 0)) = 0 & I(f(t_2, 0)) = 2 & I(f(t_3, 0)) = 1 \\ I(f(t_0, 1)) = 1 & I(f(t_1, 1)) = 3 & I(f(t_2, 1)) = 2 & I(f(t_3, 1)) = 1 \end{array}$$

令

$$\begin{array}{ll} W & = (\{0, 1, 2, 3\}, \leq) \\ w & = (0 \leq x \leq 3) \\ e & = f(b, t) \\ \varphi & = (a = s_1) \\ \psi & = (a = s_2) \\ \zeta & = (a = s_1) \end{array}$$

需要

$$\begin{array}{l} \varphi \Rightarrow (\psi \vee \zeta) \\ \zeta \Rightarrow w_x^e \wedge (\psi \vee \forall t \in T COND_t) \\ \forall t \in T, \{\zeta \wedge e = v\} t \{\psi \vee (\zeta \wedge e < v)\} \end{array}$$

试证:

$$\{\zeta \wedge e = v\} t_0 \{\psi \vee (\zeta \wedge e < v)\} .$$

需要

$$\begin{array}{l} ((a = s_1 \wedge f(b, t) = v) \wedge b = t_1 \wedge (y = 0 \vee t = 1)) \\ \rightarrow (a = s_2 \vee (a = s_1 \wedge f(t_2, t) < v)) \end{array}$$

$$\begin{array}{l} ((a = s_1 \wedge f(t_1, t) = v) \wedge b = t_1 \wedge (y = 0 \vee t = 1)) \\ \rightarrow (a = s_2 \vee (a = s_1 \wedge f(t_2, t) < v)) \end{array}$$

令

$$\begin{array}{ll} W & = (\{0, 1, 2, 3\}, \leq) \\ w & = (0 \leq x \leq 3) \\ e & = f(b, t) \\ \varphi & = (a = s_1) \\ \psi & = (a = s_2) \\ \zeta & = (a = s_1 \wedge y = 1) \end{array}$$

因此  $\varphi \Rightarrow \Diamond \psi$  .